

INTT Weekly Meeting

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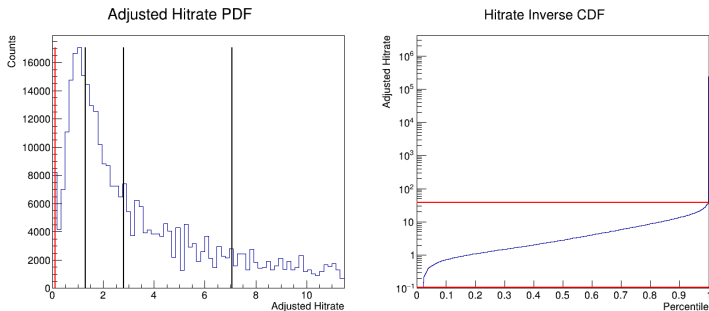
- OnlMon progress
 - Martin implemented a new API handles to retrieve all BCOs from a Packet
 - Hitmap and hitrates options are now properly normalized
 - Changes were applied yesterday, but we haven't had beam to check
- Revisiting hot channel classification
 - Using survey geometry to normalize by positions
 - Identify cold channel cutoff with local minimum in hitrate PDF
 - Identify hot channel cutoff when the tail reaches the aforementioned local minimum
- QA
 - Plots produced from calibration modules can be used for QA database
 - Calibration modules themselves can have QA functionality
 - Eventual automation of database modification with scripts and SQL
- HV SQL table
 - In the daq database of opc0
 - Logging MPOD values since April 24

- Mentioned software change I have been meaning to implement for some time
- Ultimately fell through—Chris thought it was too drastic
- Mostly agree with his decision and will revisit this after physics running (October)

- psql table in the daq database
- accessible on opc0
- table is `intt_mpodlog`
 - time timestamp
 - ip varchar(32)
 - mpod_channel smallint
 - status varchar(12)
 - voltage real
 - current real

- Updated every 2 minutes due to a cron job periodically run by John Haggerty
- The cron job runs
`home/phnxrc/INTT/josephb/HV/pycodes/mpodlog.sh`
 - Wrapper shell script that configures a `venv`
 - runs a python script
 - implementation is similar to the HV GUI
 - values obtained via `snmpwalk`
- High voltage currents/status can be monitored via Grafana
 - Nicety but I think implementing a Grafana monitor is low priority
 - But it could make a nice project for a student who wants to learn Grafana/SQL

- The fundamental shape is not described by a Gaussian or Poisson
 - Peak is too far left of the median and mean
 - The hot "tail" is too fat
- Fits using a Gaussian or Poisson will fail in several ways:
 - The mean will not correspond to the peak location, but closer to the mean
 - The variance will correspond to the IQR, and fail to predict the slow decay in the tail
 - Underestimating the variance leads to masking a higher portion of the detector
 - The fit can fail to converge normally and is subject to histogram binning



keeping $>1.051\text{E-}01$ (excludes 2.1340%)
keeping $<3.841\text{E+}01$ (excludes 0.1637%)
events: 30000

Figure: Quartiles are drawn as vertical black lines on the hitrate pdf. The peak lies outside the IQR; Gaussian fits find the mean close to the median value, with sigma commensurate with the IQR.

- Hitrate normalization:

$$h' = \frac{h \|\vec{r}\|^2}{L(\hat{n} \cdot \hat{r})} \quad (1)$$

- h' is the adjusted hitrate
- h is the raw hitrate (hits per event)
- \vec{r} is the displacement vector from an origin to the strip position
- \hat{n} is the sensor unit normal
- L is strip length, either 1.6cm for Type A or 2.0cm for Type B
- This can be understood by solving for h
- A given hitrate should be proportional to the flux one expects geometrically
- Note that \vec{r} is obtained using survey geometry assuming a vertex position $\vec{0}$
- All channels can be compared in a common distribution

Get the inverse cdf, starting with a hitmap:

```
// m_hitmap is filled over many process_event calls

// get pdf up to normalization
std::map<double, double> hitrate_pdf;
for(auto const& [channel, count] : m_hitmap) {
    // count is the number of hits seen by channel
    double h = count / m_num_evts;
    double h_adjust = adjust_hitrate(channel, h);
    ++hitrate_pdf[h_adjust];
}

// get inverse cdf
double total = 0;
double NUM_CHANNELS = 8 * 14 * 26 * 128; // number of INTT channels
std::map<double, double> hitrate_inv_cdf;
for(auto const& [hitrate, count] : hitrate_pdf) {
    total += count;
    double fraction = total / NUM_CHANNELS;
    hitrate_inv_cdf[fraction] = hitrate;
}
```

Find local minima using the inverse cdf

```
double min_hitrate = 0.0, max_hitrate = 0.0; // Our cold/hot cut values
double best_min_ratio = 0.0, best_max_ratio = 0.0;
double prev_fraction = 0.0, prev_hitrate = 0.0;
for(auto const& [fraction, hitrate] : hitrate_inv_cdf) {
    // Division is safe since the smallest fraction is >0
    double ratio = (hitrate - prev_hitrate) / (fraction - prev_fraction);

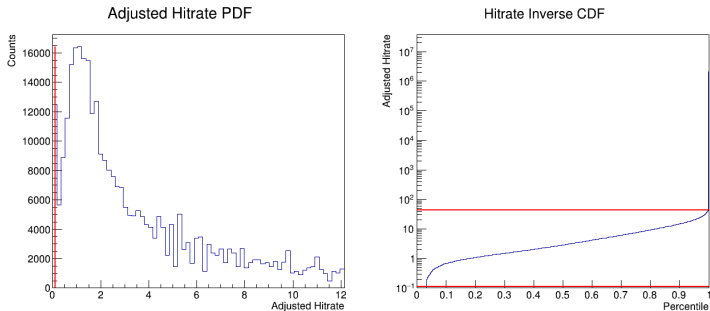
    // search for local min left of median
    if(best_min_ratio < ratio && fraction < 0.5) {
        best_min_ratio = ratio;
        min_hitrate = hitrate;
    }

    // search for max right of median
    if(best_max_ratio < ratio && 0.5 < fraction) {
        best_max_ratio = ratio;
        max_hitrate = hitrate;
    }

    // ratio can become large as the tail becomes sparse
    // stop if we've found a best_max_ratio as large as best_min_ratio
    if(best_min_ratio < best_max_ratio) {
        break;
    }
}

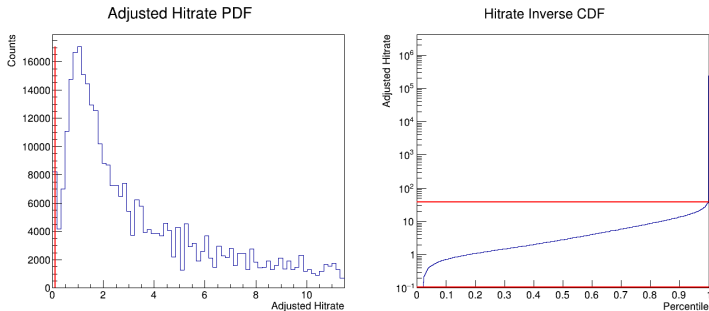
// min_hitrate is now a suitable cold cut
// max_hitrate is now a suitable hot cut
```

- Effectively, we are finding a local minimum in the pdf
- We do this by finding maximum values of the slope of the inverse cdf
- Since the pdf tail can become arbitrarily sparse, we match the value found on the left
- The algorithm is very stable and can be used for pp and $AuAu$ runs without modification
- Pooling of all channels increases statistics, and reduces the number of events needed for convergence
 - pp converge in about 30k events or less
 - $AuAu$ can converge in as few as 1k events
- Typically
 - 3%-4% of the detector is cold
 - 0.1%-0.3% of the detector is hot
 - but this is just the trend over the runs I analyzed while designing and debugging
- Examples follow on subsequent slides



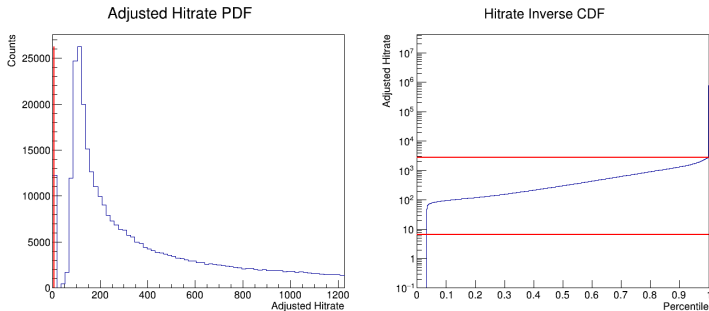
keeping $>1.102\text{E-}01$ (excludes 3.2345%)
keeping $<4.379\text{E+}01$ (excludes 0.1701%)
events: 30000

Figure: Run 41021 (pp)



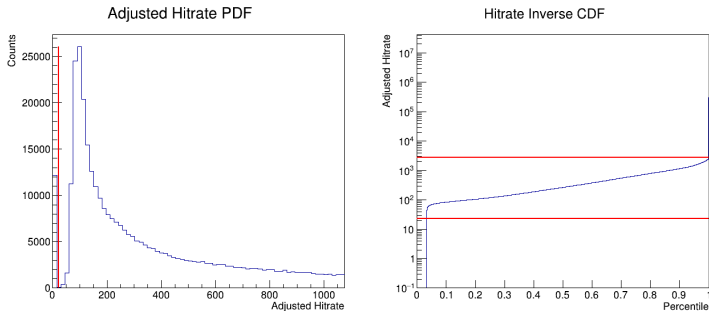
keeping $>1.051\text{E-}01$ (excludes 2.1340%)
keeping $<3.841\text{E+}01$ (excludes 0.1637%)
events: 30000

Figure: Run 42212 (pp)



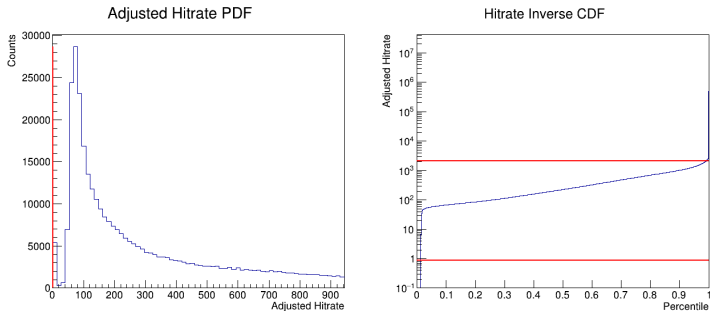
keeping $>6.601\text{E}+00$ (excludes 3.2830%)
keeping $<2.783\text{E}+03$ (excludes 0.2723%)
events: 5000

Figure: Run 20444 (*AuAu*)



keeping $>2.313E+01$ (excludes 3.2817%)
keeping $<2.699E+03$ (excludes 0.0612%)
events: 5000

Figure: Run 20446 ($AuAu$)



keeping $>8.771\text{E-}01$ (excludes 1.3613%)
keeping $<2.068\text{E}+03$ (excludes 0.7968%)
events: 5000

Figure: Run 20869 (*AuAu*)

- The plots for debugging are suitable for QA output
 - The same module/macro can be used
- Chris' students implemented a QA database
 - The natural conclusion of QA should be to issue SQL commands
 - Classify the run as "Good" or "Bad" based on module output
 - Issue SQL statements upon completion
- For the hitmap, an easy metric is the total portion of the detector that is excluded
 - For example, $> 5\%$ total excluded is "Bad"
 - Otherwise, it is "Good" unless other modules detect problems

- Implement QA modules for other things
 - BCO timing is another natural choice
 - Genki has many such modules, but they are not committed to coresoftware or macros (that I know of)
 - For each QA module, there should be a concrete conclusion of run goodness
- Continue to work on Online Monitoring
- Misc. INTT work as it arises